

## **Computational Biology** (BIOSC 1540)

### Lecture 05B

Sequence Alignment

Methodology

Feb 6, 2025



## Announcements

Assignments

- Assignment P01D is due Monday (Feb 10)
- Assignment P01E will be released on Saturday (Feb 8)?

**Quizzes** • Quiz 02 is on Feb 18 and will cover lectures 04A to 06B

#### **CBytes**

- CByte 02 expires on Feb 7
- CByte 03 expires on Feb 15
- CByte 04 releases on Feb 8

**Next reward:** Checkpoint Submission Feedback

**ATP until the next reward:** 1,653

## After today, you should have a better understanding of



#### Dynamic programming basics in bioinformatics

# Biological data analysis often requires comparing sequences, which can be computationally expensive

Sequence alignment finds the best way to arrange two sequences to maximize similarity.

**The challenge:** Finding the best alignment among exponentially many possibilities.

#### 

| Query          | 1              | ATGACTTTATCCATTCTAGTTGCACATGACTTGCAACGAGTAATTGGTTTTGAAAAATCAA | 60             |
|----------------|----------------|---|----------------|
| Sbjct          | 2555705        | AAAA.TCTCTAAAACG.ACC  | 2555646        |
| Query          | 61             | TTACCTTGGCATCTACCAAATGATTTGAAGCATGTTAAAAAATTATCAACTGGTCATACT  | 120            |
| Sbjct          | 2555645        | CTCTAAA   | 2555586        |
| Query          | 121            | TTAGTAATGGGTCGTAAGACATTTGAATCGATTGGTAAACCACTACCGAATCGTCGAAAT  | 180            |
| Sbjct          | 2555585        | C.TCAGATA.TTAGGT.GAA.   | 2555526        |
| Query<br>Sbjct | 181<br>2555525 | GTTGTACTTACTTCAGATACAAGTTTCAACGTAGAGGGCGTTGATGTAATTCATTC      | 237<br>2555469 |
| Query          | 238            | ATTGAAGATATTTATCAACTACCGGGCCATGTTTTTATATTTGGAGGGCAAACATTATTT  | 297            |
| Sbjct          | 2555468        | CTAA.AG.GTT.TTAAC.GAC   | 2555409        |
| Query          | 298            | GAAGAAATGATTGATAAAGTGGACGACATGTATATTACTGTTATTGAAGGTAAATTTCGT  | 357            |
| Sbjct          | 2555408        | CCC.GATTCAATAGAA  | 2555349        |
| Query          | 358            | GGTGATACGTTCTTTCCACCTTATACATTTGAAGACTGGGAAGTTGCCTCTTCAGTTGAA  | 417            |
| Sbjct          | 2555348        | ACAACACAC.AAA   | 2555289        |
| Query          | 418            | GGTAAACTAGATGAGAAAAATACAATTCCACATACCTTTCTACATTTAATTCGTAAAAAA  | 477            |
| Sbjct          | 2555288        | CAATAGACTG.GG   | 2555229        |

- Biological sequences (DNA, RNA, proteins) are long.
- Pairwise comparisons scale exponentially—naïve approaches take too long.
- We need efficient algorithms to compare sequences optimally.

## Comparing all possible sequence alignments grows exponentially with sequence length

Suppose we want to identify the optimal alignment for these two sequences

ATGTC ATGC

We could take the brute force approach of trying every single possible alignment and computing the score

| ATGTC | ATGTC | ATGTC | A-TGTC | ATGTC |
|-------|-------|-------|--------|-------|
| ATGC- | AT-GC | -ATGC | ATGC   | A-TGC |

The number of possible alignments is exponential for two sequences of length *m* and *n* 

$$N(m,n)pprox \sum_{d=0}^{\min(m,n)} rac{(m+n-d)!}{d!(m-d)!(n-d)!}$$

Two sequences of length 100 have  $2.05 imes 10^{75}$  possible alignments

## Dynamic programming (DP) finds the best alignment by breaking it into smaller subproblems

For our previous sequences, ATGTC this is the optimal alignment ATG-C

How can we know it's optimal if I don't try every possible combination?

| Why should we even compute this alignment | ATGTC |
|---|-------|
| score? We know it would be very low.      | -ATGC |

#### Instead of trying all alignments, **DP builds an optimal alignment step by step from the start.**

We can assert that the optimal final solution is the one where the first step (match of A) is optimal, then the second step (match of T) is optimal, etc.

Guarantees the best alignment without exhaustive searching.

# Dynamic programming constructs an optimal alignment by systematically building the alignment Alignments are built incrementally from smaller subproblems ATGTC ATGC 0. Start with an "empty" alignment and define scoring scheme

 Match: +1
 Mismatch: -1
 Gap: -2

**1.** Which move would give me the highest score for the first position?

Insert a gap?orAlignment match(This would be a sequence<br/>match or mismatch.)Let's check the first characters:Aand

This alignment match would be a sequence match; thus, our optimal alignment should start with this



**2.** Which move would give me the highest score for the second position?

Insert a gap?orAlignment matchLet's check the characters:TandT

Alignment match would be best

Next optimal alignment: A T

Sequences: ATGTC ATGCCurrent optimal alignment:A TMatch: +1Score:2A TMismatch: -1Gap: -2

**3.** Which move would give me the highest score for the third position?



Alignment match would be best

Next optimal alignment: A T G A T G



**4.** Which move would give me the highest score for the fourth position?

| Insert a gap?              | or | Alignmer | nt match |
|----------------------------|----|----------|----------|
| et's check the characters: | т  | and      | С        |

At first glance, it would seem that a sequence mismatch would be best because it would only decrease our score by one instead of two

Next optimal alignment?

ATGT ATGC However, we would need to consider what happens later

After today, you should have a better understanding of



#### Dynamic programming basics in bioinformatics

**Scoring matrix** 

# Dynamic programming constructs an optimal alignment by systematically filling a scoring matrix

- The matrix ensures that each alignment decision is optimal up to that point.
- The value in each cell reflects the best possible score for aligning the prefixes of two sequences up to that point.
- Computed using **previous scores** from adjacent cells (left, top, diagonal).
- The final score in the matrix represents the **best possible alignment score** for the entire sequences.

|   |   | 0  | 1  | 2  | 3  | 4  | 5  |
|---|---|----|----|----|----|----|----|
|   | D |    | Α  | A  | Т  | Т  | С  |
| 0 |   | 0  | -1 | -2 | -3 | -4 | -5 |
| 1 | Α | -1 | 1  | 0  | -1 | -2 | -3 |
| 2 | Т | -2 | 0  | 0  | 1  | 0  | -1 |
| 3 | T | -3 | -1 | -1 | 1  | 2  | 1  |
| 4 | A | -4 | -2 | 0  | 0  | 1  | 1  |
| 5 | С | -5 | -3 | -1 | -1 | 0  | 2  |

After today, you should have a better understanding of

#### Needleman-Wunsch algorithm (global alignment)

Let's align two sequences: AATTC ATTAC

First, enter zero in our first coordinate (0, 0)

We need to fill in each cell by moving from other cells starting from (0, 0)

Each move "uses" a nucleotide from a row, column, or both

Moving right or down uses a gap and you add the penalty to previous score



Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1



(Disclaimer: these values are not correct for the final matrix.)

**ATTAC** 

**AATTC** 

Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1



The last cell in our scoring matrix represents our final score of this alignment

Let's align two sequences:



Alignment score: -5

**ATTAC** 

**AATTC** 

Let's align two sequences:

Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1



Let's align two sequences: **AATTC** 

To fill in other cells, we need to find the best move (highest score) from **earlier, adjacent cells** 

Let's figure out this score

| Option 1         | Option 2            | Option 3            |  |  |
|------------------|---------------------|---------------------|--|--|
| A A              | Α -                 | - A                 |  |  |
| Match (+1)       | Gap (-1)            | Gap (-1)            |  |  |
| 0 + 1 = <b>1</b> | -1 + -1 = <b>-2</b> | -1 + -1 = <b>-2</b> |  |  |

ATTAC

Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1



AATTC

Let's align two sequences:

Mismatch (-1)

Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1

|                 |                 |                 |   |   | 0    | 1  | 2  | 3  | 4  | 5  |
|-----------------|-----------------|-----------------|---|---|------|----|----|----|----|----|
| <b>Option 1</b> | <b>Option 2</b> | <b>Option 3</b> |   | D |      | Α  | Α  | Т  | Т  | С  |
| <u>л</u> т      | Δ               | -               | 0 |   | 0    | -1 | -2 | -3 | -4 | -5 |
| A               | A -             |                 | 1 | Α | -1   | 1  |    |    |    |    |
| lismatch (-1)   | Gap (-1)        | Gap (-1)        | 2 | Т | -2 - | 0  |    |    |    |    |
|                 |                 |                 | 3 | Т | -3   |    | -  |    |    |    |
| -1 + -1 = -2    | -2 + -1 = -3    | 1 + -1 = 0      | 4 | Α | -4   |    |    |    |    |    |
|                 |                 |                 | 5 | С | -5   |    |    |    |    |    |

**ATTAC** 

Let's align two sequences: **AATTC** 

Repeat until we fill the matrix

ATTAC

Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1

The last number represents the best possible alignment score



We get the alignment by **tracing back** our moves to (0, 0) from our best score

Starting from the bottom left, what is the last move we made to get this score?



#### This is the last part of our alignment



Repeat for the next one



This is the second to last part of our alignment

## There can be multiple optimal alingments





# Global alignment compares sequences in their entirety

Global alignment aligns sequences **from start to end** 

Key characteristics:

- 1. Attempts to align every residue in both sequences
- 2. Introduces gaps as necessary to maintain end-to-end alignment
- 3. Optimizes the overall alignment score for the entire sequences

Guarantees finding the optimal global alignment between two sequences

## After today, you should have a better understanding of

#### Smith-Waterman algorithm (local alignment)

## Local alignment identifies best matching subsequences

Focuses on finding regions of high similarity within sequences

- Does not require aligning entire sequences end-to-end
- Allows for identification of conserved regions or domains

#### **Key characteristics:**

- Aligns subsections of sequences
- Ignores poorly matching regions
- Can find multiple areas of similarity in a single comparison

## Smith-Waterman

We have a few algorithm changes

Zero is the lowest score (i.e., if negative, make it zero) Start alignment at highest cell Stop aligning when you encounter a zero

|   |   | 0 | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|---|---|
|   | D |   | Α | Α | Т | Т | С |
| 0 |   | 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | Α | 0 | 1 | 1 | 0 | 0 | 0 |
| 2 | Т | 0 | 0 | 0 | 2 | 1 | 0 |
| 3 | Т | 0 | 0 | 0 | 1 | 3 | 2 |
| 4 | A | 0 | 1 | 1 | 0 | 2 | 2 |
| 5 | С | 0 | 0 | 0 | 0 | 1 | 3 |

Scoring scheme

- Match: +1
- Mismatch: -1
- Gap: -1



## Smith-Waterman differs from Needleman-Wunsch in key aspects

#### Matrix initialization:

- Needleman-Wunsch: The first row and column are filled with gap penalties
- Smith-Waterman: First row and column filled with zeros

#### Scoring system:

- Needleman-Wunsch: Allows negative scores
- Smith-Waterman: Negative scores are set to zero

#### Traceback:

- Needleman-Wunsch: Starts from the bottom-right cell
- Smith-Waterman: Starts from each highest-scoring cell in the matrix

## After today, you should have a better understanding of

#### Python

## Before the next class, you should

#### Lecture 05B:

Sequence alignment -Methodology

#### Lecture 06A:

Read Mapping -Foundations



• Start P01D (due Feb 10)